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Institute for Research and Applications of Fuzzy Modeling

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Jiří Močkoř

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University of Ostrava
Institute for Research and Applications of Fuzzy Modeling
30. dubna 22, 701 03 Ostrava 1, Czech Republic

tel.: +420-59-6160234 fax: +420-59-6120 478
e-mail: mockor@osu.cz

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Jiří Močkoř

University of Ostrava

Institute for Research and Applications of Fuzzy Modeling,

Department of Mathematics

30. dubna 22, 701 03 Ostrava 1, Czech Republic *

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1 Introduction

In a fuzzy set theory there are several categories which play important role in fuzzy logic interpretation. Two of these categories of fuzzy sets over MV -algebra $\Omega = (L, \otimes, \rightarrow)$ will be investigated in the paper. The first one is a category $\mathbf{Set}(\Omega)$ with objects (A, δ) where A is a set and $\delta : A \times A \rightarrow \Omega$ is a similarity relation such that

- (a) $(\forall x \in A) \quad \delta(x, x) = 1,$
- (b) $(\forall x, y \in A) \quad \delta(x, y) = \delta(y, x),$
- (c) $(\forall x, y, z \in A) \quad \delta(x, y) \otimes \delta(y, z) \leq \delta(x, z).$

A morphism $f : (A, \delta) \rightarrow (B, \gamma)$ in $\mathbf{Set}(\Omega)$ is a map $f : A \times B \rightarrow \Omega$ satisfying the following conditions

- (a) $(\forall x, z \in A)(\forall y \in B) \quad \delta(x, z) \otimes f(x, y) \leq f(z, y),$
- (b) $(\forall x \in A)(\forall y, z \in B) \quad \gamma(y, z) \otimes f(x, y) \leq f(x, z),$
- (c) $(\forall x \in A)(\forall y, z \in B) \quad f(x, y) \otimes f(x, z) \leq \gamma(y, z),$
- (d) $(\forall x \in A) \quad 1 = \bigvee \{f(x, y) : y \in B\}.$

If $f : \mathbf{A} \rightarrow \mathbf{B}$ and $g : \mathbf{B} \rightarrow \mathbf{C}$ are two morphisms then their composition is a function $g \circ f : A \times C \rightarrow \Omega$ such that

$$g \circ f(x, z) = \bigvee_{y \in B} (f(x, y) \otimes g(y, z)).$$

The other category $\mathbf{SetF}(\Omega)$ will have the same objects as the category $\mathbf{Set}(\Omega)$. A morphism $f : (A, \delta) \rightarrow (B, \gamma)$ in $\mathbf{SetF}(\Omega)$ is a map $f : A \rightarrow B$ such that $(\forall x, y \in A) \quad \gamma(f(x), f(y)) \geq \delta(x, y)$.

In [9],[10] we investigated some principal properties of a category $\mathbf{SetF}(\Omega)$ and we proved that the *extensional subobjects* of objects (A, δ) in this category $\mathbf{SetF}(\Omega)$ can be identified with some *characteristic morphism* $(A, \delta) \rightarrow (\Omega^*, \mu)$. Namely we proved that if $S : \mathbf{SetF}(\Omega) \rightarrow \mathbf{Set}$ is a functor such that $S(A, \delta) = \{s : s \text{ is an extensional subobject of } (A, \delta)\}$ then there exists a natural isomorphism

$$\zeta : S(-) \rightarrow \mathbf{Hom}_{\mathbf{SetF}(\Omega)}(-, \Omega^*).$$

This classification property, which is one of the most important properties of a topos category are frequently used for interpretation of formulas of fuzzy logic in a category $\mathbf{SetF}(\Omega)$ in such a way that interpretation of a fuzzy logic formula is defined as a special extensional subobject of some object (A, δ) .

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Hence, it seems natural that extensional subobjects of objects in a category $\mathbf{SetF}(\Omega)$ play important role for further investigation of that category.

In this paper we are interested in the following problem related to extensional subobjects. For any object (A, δ) of a category $\mathbf{SetF}(\Omega)$ we can define a set $\Omega^{(A, \delta)}$ of all (or some special, respectively) extensional subobjects of (A, δ) . This set can be transformed (in several different ways) into an object of a category $\mathbf{SetF}(\Omega)$. In that way we obtain a full subcategory $\Omega^{\mathbf{SetF}(\Omega)}$ of a category $\mathbf{SetF}(\Omega)$ consisting of these special objects. We will be interested in conditions under which that subcategory $\Omega^{\mathbf{SetF}(\Omega)}$ is a quasi-reflective subcategory in $\mathbf{SetF}(\Omega)$. Recall that a subcategory \mathbf{L} of a category \mathbf{K} is a *quasi-reflective* subcategory in \mathbf{K} if there exists a functor $G : \mathbf{K} \rightarrow \mathbf{L}$ such that for any object $a \in \mathbf{K}$ there should exist a morphism $a @ > u_a >> G(a)$ such that for any object $b \in \mathbf{L}$ and any morphism $f : a \rightarrow b$ (in \mathbf{K}) there exists a morphism (in general non unique) $\hat{f} : G(a) \rightarrow b$ such that the diagram commutes:

$$\begin{array}{ccc} a & \xrightarrow{u_a} & G(a) \\ f \downarrow & & \downarrow \hat{f} \\ b & \underline{\underline{=}} & b. \end{array}$$

A functor G is then called a *quasi-reflector*. We will be interested in several subcategories of a category $\mathbf{SetF}(\Omega)$ consisting of various objects $(sub(A, \delta), \sigma)$, where $sub(A, \delta)$ will be a subset of the set of all extensional subobjects of (A, δ) and σ will be a similarity relation defined on that subset. We prove that all these subcategories are quasi-reflective subcategories in a category $\mathbf{SetF}(\Omega)$. We also introduce a notion of a weak singleton extensional subobject of (A, δ) in a category $\mathbf{Set}(\Omega)$ and we prove that a subcategory consisting of these subobjects is also quasi-reflective subcategory of a category $\mathbf{Set}(\Omega)$.

2 Subcategories of extensional subobjects

We show firstly a simple result which states the existence of a functor between categories $\mathbf{SetF}(\Omega)$ and $\mathbf{Set}(\Omega)$.

Lemma 2.1 *There exists a functor $F : \mathbf{SetF}(\Omega) \rightarrow \mathbf{Set}(\Omega)$.*

For $(A, \delta) \in \mathbf{SetF}(\Omega)$ we set $F(A, \delta) = (A, \delta)$ and for a morphism $f : (A, \delta) \rightarrow (B, \gamma)$ in $\mathbf{SetF}(\Omega)$ we define a map $F(f) : A \times B \rightarrow \Omega$ such that $F(f)(a, b) = \gamma(f(a), b)$ for any $a \in A, b \in B$. Then $F(f)$ is a morphism in $\mathbf{Set}(\Omega)$. In fact, we have for example

$$\begin{aligned} F(f)(a, b) \otimes \delta(a, a') &= \gamma(f(a), b) \otimes \delta(a, a') \leq \\ &\leq \gamma(f(a), b) \otimes \gamma(f(a), f(a')) \leq \gamma(f(a'), b) = F(f)(a', b). \end{aligned}$$

Recall that an *extensional subobject* of (A, δ) in a category $\mathbf{SetF}(\Omega)$ is a map $s : A \rightarrow \Omega$ such that

$$s(x) \otimes \delta(x, y) \leq s(y).$$

An extensional subobject can be defined in a category $\mathbf{Set}(\Omega)$ as well. In fact, it is clear that $(\Omega, \leftrightarrow)$ is an object in $\mathbf{SetF}(\Omega)$, where $\alpha \leftrightarrow \beta = (\alpha \rightarrow \beta) \wedge (\beta \rightarrow \alpha)$. Then s is an extensional subobject of (A, δ) in a category $\mathbf{SetF}(\Omega)$ if $s : (A, \delta) \rightarrow (\Omega, \leftrightarrow)$ is a morphism in $\mathbf{SetF}(\Omega)$. Analogously s will be called an *extensional subobject* of (A, δ) in $\mathbf{Set}(\Omega)$ if $s : (A, \delta) \rightarrow (\Omega, \leftrightarrow)$ is a morphism in $\mathbf{Set}(\Omega)$, i.e.

- (a) $(\forall a, a' \in A)(\alpha \in \Omega) \quad s(a, \alpha) \otimes \delta(a, a') \leq s(a', \alpha),$
- (b) $(\forall a \in A)(\forall \beta, \alpha \in \Omega) \quad s(a, \alpha) \otimes (\alpha \leftrightarrow \beta) \leq s(a, \beta),$
- (c) $(\forall a \in A)(\alpha, \beta \in \Omega) \quad s(a, \alpha) \otimes s(a, \beta) \leq \alpha \leftrightarrow \beta,$
- (d) $(\forall x \in A) \quad 1 = \bigvee_{\alpha \in \Omega} s(x, \alpha).$

An extensional subobject s of (A, δ) in $\mathbf{SetF}(\Omega)$ is called *normal* if $\bigvee_{a \in A} s(a) = 1$. Then an extensional subobject t of (A, δ) in $\mathbf{Set}(\Omega)$ is called *normal* if $\bigvee_{a \in A} t(a, 1) = 1$. Moreover, we say that t is a *weak singleton* of (A, δ) in $\mathbf{Set}(\Omega)$ if t is a normal extensional subobject of (A, δ) in $\mathbf{Set}(\Omega)$ and $t(a, 1) \otimes t(b, 1) \leq$

$\delta(a, b)$ for all $a, b \in A$. A special normal extensional subobject in $\mathbf{SetF}(\Omega)$ is called singleton. Recall that an extensional subobject $s : A \rightarrow \Omega$ of (A, δ) (in a category $\mathbf{SetF}(\Omega)$) is a *singleton* if it satisfies the condition

$$(\forall x, y \in A) \quad s(x) \otimes s(y) \leq \delta(x, y).$$

It is clear that a map $\{a\} = \delta(a, -) : A \rightarrow \Omega$ is an example of a singleton for any $a \in A$. On the other hand an object (A, δ) is called *complete* if for any singleton s of (A, δ) there exists $a \in A$ such that $s = \delta(a, -)$.

Let (A, δ) be an object of $\mathbf{Set}(\Omega)$ (or $\mathbf{SetF}(\Omega)$). We introduce the following notation.

$$\begin{aligned} \Omega^{(A, \delta)} &= \{s : s \text{ is an extensional subobject of } (A, \delta) \text{ in } \mathbf{SetF}(\Omega)\}, \\ \Omega_1^{(A, \delta)} &= \{s : s \text{ is an extensional and normal subobject of } (A, \delta) \text{ in } \mathbf{SetF}(\Omega)\}, \\ \text{w-singl}(A, \delta) &= \{s : s \text{ is a weak singleton of } (A, \delta) \text{ in } \mathbf{Set}(\Omega)\}, \\ \text{singl}(A, \delta) &= \{s \in \Omega_1^{(A, \delta)} : s \text{ is singleton of } (A, \delta) \text{ in } \mathbf{SetF}(\Omega)\}. \end{aligned}$$

All the previous sets can be transformed into objects of categories $\mathbf{Set}(\Omega)$ and $\mathbf{SetF}(\Omega)$, respectively. In fact, for any object (A, δ) and for any $s, t \in \Omega^{(A, \delta)}$, $p, q \in \text{w-singl}(A, \delta)$ we set

$$\begin{aligned} \sigma(s, t) &= \sigma_{(A, \delta)}(s, t) = \bigwedge_{x \in A} s(x) \leftrightarrow t(x), \\ \tau(s, t) &= \tau_{(A, \delta)}(s, t) = \begin{cases} \bigvee_{x \in A} s(x) \otimes t(x), & \text{if } s \neq t \\ 1, & \text{if } s = t \end{cases}, \\ \rho(p, q) &= \rho_{(A, \delta)}(p, q) = \bigwedge_{x \in A} p(x, 1) \leftrightarrow q(x, 1). \end{aligned}$$

Lemma 2.2 *For any object (A, δ) there exists a morphism $\hat{} : (\text{singl}(A, \delta), \sigma_{(A, \delta)}) \rightarrow (\text{w-singl}(A, \delta), \rho_{(A, \delta)})$.*

Proof. For $s \in \text{singl}(A, \delta)$ we set $\hat{s}(a, \alpha) = s(a) \leftrightarrow \alpha$. It is then clear that $\hat{s} = F(s)$ (see Lemma 1), since $s : (A, \delta) \rightarrow (\Omega, \leftrightarrow)$ is a morphism in $\mathbf{SetF}(\Omega)$. It follows that \hat{s} is an extensional (and clearly normal) subobject in $\mathbf{Set}(\Omega)$. Since s is a singleton, \hat{s} is a weak singleton. Moreover according to Lemma 5 for any $s, t \in \Omega_1^{(A, \delta)}$ we have

$$\begin{aligned} \rho_{(A, \delta)}(\hat{s}, \hat{t}) &= \bigwedge_{a \in A} (s(a) \leftrightarrow 1) \leftrightarrow (t(a) \leftrightarrow 1) = \\ &= \bigwedge_{a \in A} s(a) \leftrightarrow t(a) = \sigma_{(A, \delta)}(s, t). \end{aligned}$$

Lemma 2.3 *For any object (A, δ) the pairs $(\Omega^{(A, \delta)}, \sigma_{(A, \delta)})$, $(\Omega^{(A, \delta)}, \tau_{(A, \delta)})$, $(\text{w-singl}(A, \delta), \rho_{(A, \delta)})$ and $(\text{singl}(A, \delta), \tau_{(A, \delta)})$, respectively are objects of a category $\mathbf{Set}(\Omega)$ (and $\mathbf{SetF}(\Omega)$, simultaneously).*

The proof of this lemma can be done by simple computation.

Lemma 2.4 *Let Ω be an MV-algebra. Let $\{a_i\}_{i \in I}$ and $\{b_i\}_{i \in I}$ be two sets of elements of Ω and let $p : I \rightarrow I$ be a bijection map.*

- (a) $\bigvee_{i \in I} a_i \leftrightarrow \bigvee_{i \in I} b_{i \circ p} \geq \bigwedge_{j \in I} (a_j \leftrightarrow b_{p(j)})$,
- (b) $\bigwedge_{i \in I} a_i \leftrightarrow \bigwedge_{i \in I} b_i \geq \bigwedge_{j \in I} (a_j \leftrightarrow b_{p(j)})$.

The proof can be done by simple computation and will be omitted.

Lemma 2.5 *Let Ω be an MV-algebra.*

- (a) $(\forall a, b, x \in \Omega) \quad (x \leftrightarrow a) \leftrightarrow (x \leftrightarrow b) \geq a \leftrightarrow b$.

$$(b) (\forall a, b, c, d \in \Omega) \quad (a \leftrightarrow c) \leftrightarrow (b \leftrightarrow d) \geq (a \leftrightarrow b) \otimes (c \leftrightarrow d).$$

(1) We have

$$\begin{aligned} (x \rightarrow a) \leftrightarrow (x \rightarrow b) &\geq (x \leftrightarrow x) \otimes (a \leftrightarrow b) = a \leftrightarrow b, \\ (a \rightarrow x) \leftrightarrow (b \rightarrow x) &\geq (a \leftrightarrow b) \otimes (x \leftrightarrow x) = a \leftrightarrow b, \end{aligned}$$

and it follows

$$\begin{aligned} (x \leftrightarrow a) \leftrightarrow (x \leftrightarrow b) &= ((x \rightarrow a) \wedge (a \rightarrow x)) \leftrightarrow ((x \rightarrow b) \wedge (b \rightarrow x)) \geq \\ &((x \rightarrow a) \leftrightarrow (x \rightarrow b)) \wedge ((a \rightarrow x) \leftrightarrow (b \rightarrow x)) \geq a \leftrightarrow b. \end{aligned}$$

(2) According to Lemma 4 we have

$$\begin{aligned} (a \leftrightarrow c) \leftrightarrow (b \leftrightarrow d) &= ((a \rightarrow c) \wedge (c \rightarrow a)) \leftrightarrow ((b \rightarrow d) \wedge (d \rightarrow b)) \geq \\ &((a \rightarrow c) \leftrightarrow (b \rightarrow d)) \wedge ((c \rightarrow a) \leftrightarrow (d \rightarrow b)) \geq \\ &((a \leftrightarrow b) \otimes (c \leftrightarrow d)) \wedge ((c \leftrightarrow d) \otimes (a \leftrightarrow b)) = (a \leftrightarrow b) \otimes (c \leftrightarrow d). \end{aligned}$$

We will consider the following subcategories of categories $\mathbf{Set}(\Omega)$ and $\mathbf{SetF}(\Omega)$, respectively.

- (a) A full subcategory $\mathbf{SetF}(\Omega)_{comp} \hookrightarrow \mathbf{SetF}(\Omega)$ consisting of complete objects of a category $\mathbf{SetF}(\Omega)$,
- (b) A full subcategory $\Omega_{\leftrightarrow}^{\mathbf{SetF}(\Omega)} \hookrightarrow \mathbf{SetF}(\Omega)$ with objects $(\Omega^{(A,\delta)}, \sigma_{(A,\delta)})$ for any object (A, δ) ,
- (c) A full subcategory $\Omega_{\otimes}^{\mathbf{SetF}(\Omega)} \hookrightarrow \mathbf{SetF}(\Omega)$ with objects $(\Omega^{(A,\delta)}, \tau_{(A,\delta)})$ for any object (A, δ) ,
- (d) A full subcategory $\Omega_{1, \otimes}^{\mathbf{Set}(\Omega)} \hookrightarrow \mathbf{Set}(\Omega)$ with objects $(\Omega_1^{(A,\tau)}, \sigma_{(A,\tau)})$ for any object (A, δ) ,
- (e) A full subcategory $\Omega^{\mathbf{Set}(\Omega)} \hookrightarrow \mathbf{Set}(\Omega)$ with objects $(\text{w-singl}(A, \delta), \rho_{(A,\delta)})$ for any object (A, δ) .

Theorem 2.1 *There is a functor $C : \mathbf{SetF}(\Omega) \rightarrow \mathbf{SetF}(\Omega)_{comp}$ which is a quasi-reflector.*

Let (A, δ) be an object in $\mathbf{SetF}(\Omega)$. We show firstly that $(\text{singl}(A, \delta), \tau)$ is a complete object. Let S be a singleton in $(\text{singl}(A, \delta), \tau_{(A,\delta)})$. Then we define a map $e_S : A \rightarrow \Omega$ such that

$$e_S(x) = \bigvee_{t \in \text{singl}(A, \delta)} t(x) \otimes S(t).$$

We show that $e_S \in \text{singl}(A, \delta)$. It can be easy to see that e_S is a normal extensional subobject. Moreover, we have

$$\begin{aligned} e_S(x) \otimes e_S(y) &= \bigvee_{t, p \in \text{singl}(A, \delta)} t(x) \otimes p(y) \otimes S(t) \otimes S(p) \leq \\ \bigvee_{t, p \in \text{singl}(A, \delta)} t(x) \otimes p(y) \otimes \tau(t, p) &= \bigvee_{t, p \in \text{singl}(A, \delta)} t(x) \otimes \left(\bigvee_{a \in A} t(a) \otimes p(a) \otimes p(y) \right) \leq \\ \bigvee_{t, p \in \text{singl}(A, \delta)} t(x) \otimes \left(\bigvee_{a \in A} t(a) \otimes \delta(a, y) \right) &\leq \bigvee_{t \in \text{singl}(A, \delta)} t(x) \otimes t(y) \leq \delta(x, y). \end{aligned}$$

Then $S = \{e_S\}$. In fact, let $s \in \text{singl}(A, \delta)$, then we have

$$\begin{aligned} \{e_S\}(s) &= \tau_{(A,\delta)}(e_S, s) = \bigvee_{x \in A} \bigvee_{t \in \text{singl}(A, \delta)} t(x) \otimes S(t) \otimes s(x) = \\ &\bigvee_{t \in \text{singl}(A, \delta)} \left(\bigvee_{x \in A} t(x) \otimes s(x) \right) \otimes S(t) = \\ \bigvee_{t \in \text{singl}(A, \delta)} \tau_{(A,\delta)}(s, t) \otimes S(t) &\geq \tau_{(A,\delta)}(s, s) \otimes S(s) = S(s), \end{aligned}$$

and on the other hand since $\tau(s, t) \otimes S(t) \leq S(s)$, we obtain that $\{e_S\}(s) = S(s)$. We define a functor

$$C : \mathbf{SetF}(\Omega) \rightarrow \mathbf{SetF}(\Omega)_{comp}$$

such that $C(A, \delta) = (\text{singl}(A, \delta), \tau_{(A, \delta)})$ and for a morphism $f : (A, \delta) \rightarrow (B, \gamma)$ in $\mathbf{SetF}(\Omega)$ we set $C(f) = \bar{f}$, where $\bar{f}(s)(b) = \bigvee_{a \in A} s(a) \otimes \gamma(f(a), b)$. Then \bar{f} is a morphism in $\mathbf{SetF}(\Omega)_{comp}$. In fact it is clear that $\bar{f}(s)$ is a normal extensional subobject in (B, γ) . For any $b, c \in B$ we have

$$\begin{aligned} \bar{f}(s)(b) \otimes \bar{f}(s)(c) &= \bigvee_{x, y \in A} s(x) \otimes s(y) \otimes \gamma(f(x), b) \otimes \gamma(f(y), c) \leq \\ &\quad \bigvee_{x, y \in A} \delta(x, y) \otimes \gamma(f(x), b) \otimes \gamma(f(x), c) \leq \\ &\quad \bigvee_{x, y \in A} \gamma(f(x), f(y)) \otimes \gamma(f(x), b) \otimes \gamma(f(y), c) \leq \\ &\quad \bigvee_{x, y \in A} \gamma(b, f(y)) \otimes \gamma(f(y), c) \leq \gamma(b, c). \end{aligned}$$

Hence $\bar{f}(s)$ is a singleton in (B, γ) . We show further that $\bar{f} : (\text{singl}(A, \delta), \tau_{(A, \delta)}) \rightarrow (\text{singl}(B, \gamma), \tau_{(B, \gamma)})$ is a morphism in $\mathbf{SetF}(\Omega)$. In fact, we have

$$\begin{aligned} \tau_{(B, \gamma)}(\bar{f}(s), \bar{f}(t)) &= \bigvee_{b \in B} \bar{f}(s)(b) \otimes \bar{f}(t)(b) = \\ &= \bigvee_{b \in B} \left(\bigvee_{x \in A} s(x) \otimes \gamma(f(x), b) \right) \otimes \left(\bigvee_{y \in A} t(y) \otimes \gamma(f(y), b) \right) \geq \\ &= \bigvee_{b \in B} \bigvee_{x \in A} s(x) \otimes \gamma(f(x), b) \otimes t(x) \otimes \gamma(f(x), b) \geq \\ &= \bigvee_{x \in A} s(x) \otimes \gamma(f(x), f(x)) \otimes t(x) \otimes \gamma(f(x), f(x)) = \\ &= \bigvee_{x \in A} s(x) \otimes t(x) = \tau_{(A, \delta)}(s, t). \end{aligned}$$

Let us now consider the singleton map

$$(A, \delta) \rightarrow C(A, \delta) = (\text{singl}(A, \delta), \tau_{(A, \delta)}).$$

We show that C is a quasi-reflector. Since $\tau(\{x\}, \{y\}) = \delta(x, y)$, it is clear that $\{-\}$ is a morphism in $\mathbf{SetF}(\Omega)$. Moreover let (B, γ) be a complete object and let $f : (A, \delta) \rightarrow (B, \gamma)$ be a morphism in $\mathbf{SetF}(\Omega)$. Then there exists a morphism $\tilde{f} : (\text{singl}(A, \delta), \tau) \rightarrow (B, \gamma)$ such that the diagram commutes:

$$\begin{array}{ccc} (A, \delta) & \xrightarrow{\{-\}} & (\text{singl}(A, \delta), \tau_{(A, \delta)}) \\ f \downarrow & & \downarrow \tilde{f} \\ (B, \gamma) & \xlongequal{\quad} & (B, \gamma) \end{array}$$

The map \tilde{f} is defined as follows. Let $s \in \text{singl}(A, \delta)$. Then $C(f)(s) = \bar{f}(s) : B \rightarrow \Omega$ is a singleton in (B, γ) . Since (B, γ) is complete, there exists the unique element $b \in B$ such that $\bar{f}(s) = \{b\}$. We set $\tilde{f}(s) = b$. We show that \tilde{f} is a morphism in $\mathbf{SetF}(\Omega)$. In fact, let $\tilde{f}(s) = b, \tilde{f}(t) = c$. Then we have $\gamma(\tilde{f}(s), \tilde{f}(t)) = \gamma(b, c) = \tau_{(B, \gamma)}(\{b\}, \{c\}) = \tau_{(B, \gamma)}(\bar{f}(s), \bar{f}(t)) \geq \tau_{(A, \delta)}(s, t)$ since \bar{f} is a morphism in $\mathbf{SetF}(\Omega)$. We show that the above mentioned diagram commutes. In fact, let $a \in A$, then we have $\tilde{f}(\{a\}) = b$, where $\bar{f}(\{a\}) = \{b\}$. But we have $\bar{f}(\{a\})(y) = \bigvee_{x \in A} \{a\}(x) \otimes \gamma(f(x), y) = \bigvee_{x \in A} \delta(a, x) \otimes \gamma(f(x), y) \geq \delta(a, a) \otimes \gamma(f(a), y) = \gamma(f(a), y) = \{f(a)\}(y)$. On the other hand we have $\bar{f}(\{a\})(y) \leq \bigvee_{x \in A} \gamma(f(a), f(x)) \otimes \gamma(f(x), y) \leq \bigvee_{x \in A} \gamma(f(a), y) = \gamma(f(a), y) = \{f(a)\}(y)$. Hence we have $\{b\} = \bar{f}(\{a\}) = \{f(a)\}$ and it follows that $b = f(a)$.

Theorem 2.2 *There exists a functor $D : \mathbf{SetF}(\Omega) \rightarrow \Omega_{\leftrightarrow}^{\mathbf{SetF}(\Omega)}$ which is a quasi-reflector.*

Let (A, δ) be an object of $\mathbf{SetF}(\Omega)$ and let $f : (A, \delta) \rightarrow (B, \gamma)$ be a morphism in $\mathbf{SetF}(\Omega)$. We define a functor $D : \mathbf{SetF}(\Omega) \rightarrow \Omega_{\leftrightarrow}^{\mathbf{SetF}(\Omega)}$ such that

$$D(A, \delta) = (\Omega^{(A, \delta)}, \sigma_{(A, \delta)}), \quad D(f) : D(A, \delta) \rightarrow D(B, \gamma),$$

$$(\forall s \in \Omega^{(A, \delta)})(\forall b \in B) \quad D(f)(s)(b) = \bigvee_{x \in A} s(x) \otimes \gamma(b, f(x)).$$

It is clear that this definition is correct.

Now let (A, δ) be an object in $\mathbf{SetF}(\Omega)$. We consider a map

$$(A, \delta) \rightarrow D(A, \delta) = (\Omega^{(A, \delta)}, \sigma_{(A, \delta)}).$$

We show that this map is a morphism in $\mathbf{SetF}(\Omega)$. In fact, for $x, y \in A$ we have

$$\sigma_{(A, \delta)}(\{x\}, \{y\}) = \bigwedge_{a \in A} \delta(a, x) \leftrightarrow \delta(a, y) \geq$$

$$\bigwedge_{a \in A} (\delta(a, x) \rightarrow \delta(x, y) \otimes \delta(a, x)) \wedge (\delta(a, y) \rightarrow \delta(a, y) \otimes \delta(x, y)) \geq \delta(x, y).$$

On the other hand we have $\sigma(\{x\}, \{y\}) \leq \delta(x, y)$ and it follows that $\sigma(\{x\}, \{y\}) = \delta(x, y)$.

Finally, let $f : (A, \delta) \rightarrow (\Omega^{(B, \gamma)}, \sigma_{(B, \gamma)})$ be a morphism in $\mathbf{SetF}(\Omega)$. Then there exists a morphism \widehat{f} such that the following diagram commutes.

$$\begin{array}{ccc} (A, \delta) & \xrightarrow{\{-\}} & D(A, \delta) \\ f \downarrow & & \downarrow \widehat{f} \\ (\Omega^{(B, \gamma)}, \sigma_{(B, \gamma)}) & \xlongequal{\quad} & (\Omega^{(B, \gamma)}, \sigma_{(B, \gamma)}) \end{array}$$

The morphism $\widehat{f} : \Omega^{(A, \delta)} \rightarrow \Omega^{(B, \gamma)}$ is defined as follows.

$$(\forall s \in \Omega^{(A, \delta)})(\forall b \in B) \quad \widehat{f}(s)(b) = \bigvee_{x \in A} f(x)(b) \otimes \sigma_{(A, \delta)}(\{x\}, s).$$

This definition is correct. In fact, we show firstly that $\widehat{f}(s) \in \Omega^{(B, \gamma)}$. Let $b, c \in B$, then we have

$$\widehat{f}(s)(b) \otimes \gamma(b, c) = \bigvee_{x \in A} f(x)(b) \otimes \sigma_{(A, \delta)}(\{x\}, s) \otimes \gamma(b, c) \leq$$

$$\bigvee_{x \in A} f(x)(c) \otimes \sigma_{(A, \delta)}(\{x\}, s) = \widehat{f}(s)(c).$$

Further, \widehat{f} is a morphism in $\mathbf{SetF}(\Omega)$. In fact, we have

$$\sigma_{(B, \gamma)}(\widehat{f}(s), \widehat{f}(t)) = \bigwedge_{b \in B} \widehat{f}(s)(b) \leftrightarrow \widehat{f}(t)(b).$$

According to Lemma 4 and Lemma 5 we have

$$\begin{aligned}
& \widehat{f}(s)(b) \leftrightarrow \widehat{f}(t)(b) = \\
& (\bigvee_{a \in A} f(a)(b) \otimes \sigma_{(A,\delta)}(\{a\}, s)) \leftrightarrow (\bigvee_{c \in A} f(c)(b) \otimes \sigma_{(A,\delta)}(\{c\}, t)) \geq \\
& \bigwedge_{x \in A} (f(x)(b) \otimes \sigma_{(A,\delta)}(\{x\}, s) \leftrightarrow f(x)(b) \otimes \sigma_{(A,\delta)}(\{x\}, t)) \geq \\
& \bigwedge_{a \in A} \sigma_{(A,\delta)}(\{a\}, s) \leftrightarrow \sigma_{(A,\delta)}(\{a\}, t) = \\
& \bigwedge_{a \in A} ((\bigwedge_{x \in A} \delta(a, x) \leftrightarrow s(x)) \leftrightarrow (\bigwedge_{y \in A} \delta(a, y) \leftrightarrow t(y))) \geq \\
& \bigwedge_{a \in A} \bigwedge_{x \in A} (\delta(a, x) \leftrightarrow s(x)) \leftrightarrow (\delta(a, x) \leftrightarrow t(x)) \geq \\
& \bigwedge_{a \in A} \bigwedge_{x \in A} (\delta(a, x) \leftrightarrow \delta(a, x)) \otimes (s(x) \leftrightarrow t(x)) = \\
& \bigwedge_{x \in A} s(x) \leftrightarrow t(x) = \sigma_{(A,\delta)}(s, t).
\end{aligned}$$

It follows that $\sigma_{(B,\gamma)}(\widehat{f}(s), \widehat{f}(t)) \geq \sigma_{(A,\delta)}(s, t)$. We show that the diagram commutes. In fact, let $a \in A, b \in B$. Then

$$\begin{aligned}
& \widehat{f}(\{a\})(b) = \bigvee_{x \in A} f(x)(b) \otimes \sigma_{(A,\delta)}(\{x\}, \{a\}) = \\
& \bigvee_{x \in A} f(x)(b) \otimes \delta(x, a) \leq \bigvee_{x \in A} f(x)(b) \otimes \sigma_{(B,\gamma)}(f(x), f(a)) = \\
& \bigvee_{x \in A} f(x)(b) \otimes (\bigwedge_{y \in B} f(x)(y) \leftrightarrow f(a)(y)) = \\
& \bigvee_{x \in A} \bigwedge_{y \in B} f(x)(b) \otimes (f(x)(y) \leftrightarrow f(a)(y)) \leq \bigvee_{x \in A} f(x)(b) \otimes (f(x)(b) \leftrightarrow f(a)(b)) \leq \\
& \bigvee_{x \in A} f(x)(b) \otimes (f(x)(b) \rightarrow f(a)(b)) \leq f(a)(b)
\end{aligned}$$

On the other hand we have

$$\widehat{f}(\{a\})(b) = \bigvee_{x \in A} f(x)(b) \otimes \delta(x, a) \geq f(a)(b).$$

Hence the diagram commutes and $D : \mathbf{SetF}(\Omega) \rightarrow \Omega^{\mathbf{SetF}(\Omega)}$ is a quasi-reflector.

It should be observed that \widehat{f} is the smallest morphism for which the above diagram commutes. In fact, let $g : D(A, \delta) \rightarrow (\Omega^{(B,\gamma)}, \sigma_{(B,\gamma)})$ be a morphism in $\mathbf{SetF}(\Omega)$ such that the diagram commutes. Then for any $s \in \Omega^{(A,\delta)}, b \in B$ we have

$$\begin{aligned}
& \widehat{f}(s)(b) = \bigvee_{a \in A} f(a)(b) \otimes \sigma_{(A,\delta)}(\{a\}, s) \leq \\
& \bigvee_{a \in A} f(a)(b) \otimes \sigma_{(B,\gamma)}(g(\{a\}), g(s)) = \bigvee_{a \in A} f(a)(b) \otimes \sigma_{(B,\gamma)}(f(a), g(s)) \leq \\
& \bigvee_{a \in A} f(a)(b) \otimes \bigwedge_{x \in B} f(a)(x) \rightarrow g(s)(x) \leq \bigvee_{a \in A} f(a)(b) \otimes (f(a)(b) \rightarrow g(s)(b)) \leq \\
& g(s)(b).
\end{aligned}$$

Theorem 2.3 *There exists a functor $E : \mathbf{SetF}(\Omega) \rightarrow \Omega_{\otimes}^{\mathbf{SetF}(\Omega)}$ which is a quasi-reflector.*

Let (A, δ) be an object in $\mathbf{SetF}(\Omega)$ and let $f : (A, \delta) \rightarrow (B, \gamma)$ be a morphism in $\mathbf{SetF}(\Omega)$. We define a functor E such that

$$E(A, \delta) = (\Omega^{(A, \delta)}, \tau_{(A, \delta)}), \quad E(f) : E(A, \delta) \rightarrow E(B, \gamma),$$

$$(\forall s \in \Omega^{(A, \delta)})(\forall b \in B) \quad D(f)(s)(b) = \bigvee_{x \in A} s(x) \otimes \gamma(b, f(x)).$$

It is clear that $\tau_{(A, \delta)}(s, t) \leq \tau_{(B, \gamma)}(E(f)(s), E(f)(t))$ for any $s, t \in \Omega^{(A, \delta)}$ and it follows that E is defined correctly. We consider a map

$$(A, \delta) \rightarrow (\Omega^{(A, \delta)}, \tau_{(A, \delta)})$$

such that $\{-\}(a)(b) = \tau_{(A, \delta)}(\{a\}, \{b\})$. Since $\tau_{(A, \delta)}(\{a\}, \{b\}) = \delta(a, b)$, the definition is correct.

Finally, let $f : (A, \delta) \rightarrow (\Omega^{(B, \gamma)}, \tau_{(B, \gamma)})$ be a morphism in $\mathbf{SetF}(\Omega)$. Then there exists a morphism \widehat{f} such that the following diagram commutes.

$$\begin{array}{ccc} (A, \delta) & \xrightarrow{\{-\}} & E(A, \delta) \\ f \downarrow & & \downarrow \widehat{f} \\ (\Omega^{(B, \gamma)}, \tau_{(B, \gamma)}) & \xlongequal{\quad} & (\Omega^{(B, \gamma)}, \tau_{(B, \gamma)}) \end{array}$$

The morphism $\widehat{f} : E(A, \delta) \rightarrow \Omega^{(B, \gamma)}$ is defined as follows.

$$(\forall s \in \Omega^{(A, \delta)})(\forall b \in B) \quad \widehat{f}(s)(b) = \bigvee_{x \in A} f(x)(b) \otimes \tau_{(A, \delta)}(\{x\}, s) = \bigvee_{x \in A} f(x)(b) \otimes s(x).$$

We show that \widehat{f} is a morphism in $\mathbf{SetF}(\Omega)$. In fact, let $s, t \in \Omega^{(A, \delta)}$. Then we have

$$\begin{aligned} \tau_{(B, \gamma)}(\widehat{f}(s), \widehat{f}(t)) &= \bigvee_{y \in B} \bigvee_{x, z \in A} f(x)(y) \otimes f(z)(y) \otimes s(x) \otimes t(z) = \\ &= \bigvee_{x, z \in A} \left(\bigvee_{y \in B} f(x)(y) \otimes f(z)(y) \right) \otimes s(x) \otimes t(z) = \\ &= \bigvee_{x, z \in A} \tau_{(B, \gamma)}(f(x), f(z)) \otimes s(x) \otimes t(z) \geq \bigvee_{x, z \in A} \delta(x, z) \otimes s(x) \otimes t(z) \\ &\geq \bigvee_{x \in A} s(x) \otimes t(x) = \tau_{(A, \delta)}(s, t). \end{aligned}$$

The above mentioned diagram commutes. In fact, let $x \in A, b \in B$. Then we have

$$\begin{aligned} \widehat{f}(\{a\})(b) &= \bigvee_{x \in A} f(x)(b) \otimes \{a\}(x) \geq f(a)(b), \\ \widehat{f}(\{a\})(b) &= \bigvee_{x \in A} f(x)(b) \otimes \delta(a, x) \leq \\ &= \bigvee_{x \in A} f(x)(b) \otimes \tau_{(B, \gamma)}(f(x), f(a)) \leq f(a)(b). \end{aligned}$$

Hence, $E : \mathbf{SetF}(\Omega) \rightarrow \Omega_{\otimes}^{\mathbf{SetF}(\Omega)}$ is a quasi-reflector.

Theorem 2.4 *There exists a functor $G : \mathbf{Set}(\Omega) \rightarrow \Omega_{1, \otimes}^{\mathbf{Set}(\Omega)}$ which is a quasi-reflector.*

The functor G is defined such that $G(A, \delta) = (\Omega_1^{(A, \delta)}, \tau_{(A, \delta)})$ which is considered as an object of a category $\mathbf{Set}(\Omega)$. Let $f : (A, \delta) \rightarrow (B, \gamma)$ be a morphism in $\mathbf{Set}(\Omega)$. The morphism $G(f)$ will be defined later. Firstly let us define a morphism $v_{(A, \delta)} = v : (A, \delta) \rightarrow G(A, \delta)$ in $\mathbf{Set}(\Omega)$ such that $v = F(\{-\})$, where $\{-\} : (A, \delta) \rightarrow (\Omega_1^{(A, \delta)}, \tau_{(A, \delta)})$ is a singleton morphism in $\mathbf{SetF}(\Omega)$ and F is a functor from Lemma 1. This definition is correct since for any $a, b \in A$ we have $\tau_{(A, \delta)}(\{a\}, \{b\}) = \delta(a, b)$ as it can be proved simply. Then for any object (B, γ) and any morphism $f : (A, \delta) \rightarrow (\Omega_1^{(B, \gamma)}, \tau_{(B, \gamma)})$ in $\mathbf{Set}(\Omega)$ there exists a morphism \tilde{f} in $\mathbf{Set}(\Omega)$ such that the diagram commutes.

$$\begin{array}{ccc} (A, \delta) & \xrightarrow{v_{(A, \delta)}} & G(A, \delta) \\ f \downarrow & & \downarrow \tilde{f} \\ (\Omega_1^{(B, \gamma)}, \tau_{(B, \gamma)}) & \xlongequal{\quad} & (\Omega_1^{(B, \gamma)}, \tau_{(B, \gamma)}) \end{array}$$

In fact, we set

$$(\forall s \in \Omega_1^{(A, \delta)})(\forall t \in \Omega_1^{(B, \gamma)}) \quad \tilde{f}(s, t) = \bigvee_{x \in A} f(x, t) \otimes v_{(A, \delta)}(x, s).$$

Then we have

$$\begin{aligned} \bigvee_{t \in \Omega_1^{(B, \gamma)}} \tilde{f}(s, t) &= \bigvee_{t \in \Omega_1^{(B, \gamma)}} \bigvee_{x \in A} f(x, t) \otimes v(x, s) = \\ \bigvee_{x \in A} \left(\bigvee_{t \in \Omega_1^{(B, \gamma)}} f(x, t) \right) \otimes v(x, s) &= \bigvee_{x \in A} v(x, s) = \bigvee_{x \in A} \tau_{(A, \delta)}(\{x\}, s) \geq \\ & \bigvee_{x \in A} s(x) = 1 \end{aligned}$$

Further,

$$\begin{aligned} \tilde{f}(s, t) \otimes \tilde{f}(s, t') &= \bigvee_{a, b \in A} f(a, t) \otimes f(b, t') \otimes v(a, s) \otimes v(b, s) \leq \\ \bigvee_{a, b \in A} f(a, t) \otimes f(b, t') \otimes \delta(a, b) &\leq \bigvee_{b \in A} f(b, t) \otimes f(b, t') \leq \tau_{(A, \delta)}(t, t'). \end{aligned}$$

The other properties of a morphism in $\mathbf{Set}(\Omega)$ can be proved analogously.

We show that the diagram commutes. In fact, let $a \in A, t \in \Omega_1^{(B, \gamma)}$. Then we have

$$\begin{aligned} (\tilde{f} \circ v)(a, t) &= \bigvee_{s \in \Omega_1^{(A, \delta)}} v(a, s) \otimes \tilde{f}(s, t) = \\ \bigvee_{s \in \Omega_1^{(A, \delta)}} v(a, s) \otimes \left(\bigvee_{x \in A} f(x, t) \otimes v(x, s) \right) &\geq \\ v(a, \{a\}) \otimes \bigvee_{x \in A} f(x, t) \otimes v(x, \{a\}) &= \bigvee_{x \in A} f(x, t) \otimes v(x, \{a\}) \geq \\ & f(a, t), \end{aligned}$$

and on the other hand we have

$$\begin{aligned} \tilde{f} \circ v(a, t) &= \bigvee_{s \in \Omega_1^{(A, \delta)}} \bigvee_{x \in A} \tau_{(A, \delta)}(\{a\}, s) \otimes \tau_{(A, \delta)}(\{x\}, s) \otimes f(x, t) \leq \\ \bigvee_{s \in \Omega_1^{(A, \delta)}} \bigvee_{x \in A} \tau_{(A, \delta)}(\{a\}, \{x\}) \otimes f(x, t) &= \bigvee_{x \in A} \delta(a, x) \otimes f(x, t) \leq f(a, t). \end{aligned}$$

The morphism $G(f) : G(A, \delta) \rightarrow G(B, \gamma)$ will be defined such that $G(f) = \widetilde{v_{(B, \gamma)}} \circ f$. Hence, more explicitly, we have

$$G(f)(s, t) = \bigvee_{x \in A} \bigvee_{y \in B} f(x, y) \otimes v_{(B, \gamma)}(y, t) \otimes v_{(A, \delta)}(x, s).$$

Theorem 2.5 *There exists a functor $H : \mathbf{Set}(\Omega) \rightarrow \Omega^{\mathbf{Set}(\Omega)}$ which is a quasi-reflector.*

Let (A, δ) be an object in $\mathbf{Set}(\Omega)$. We set $H(A, \delta) = (\text{w-singl}(A, \delta), \rho_{(A, \delta)})$. For an element $a \in A$ we define a morphism $[a] : (A, \delta) \rightarrow (\Omega, \leftrightarrow)$ in a category $\mathbf{Set}(\Omega)$ such that $[a](x, \alpha) = F(\{a\})(x, \alpha)$, for any $x \in A, \alpha \in \Omega$, where F is a functor from Lemma 1 and $\{a\} : (A, \delta) \rightarrow (\Omega, \leftrightarrow)$ is a morphism in a category $\mathbf{SetF}(\Omega)$ such that $\{a\}(x) = \delta(a, x)$. Then $[a]$ is a normal extensional subobject of (A, δ) in a category $\mathbf{Set}(\Omega)$. Since $[a](x, 1) \otimes [a](y, 1) \leq \delta(x, y)$ we obtain that $[a]$ is a weak singleton. Moreover, since $\rho_{(A, \delta)}([a], [b]) = \delta(a, b)$ we obtain that $[-] : (A, \delta) \rightarrow (\text{w-singl}(A, \delta), \rho_{(A, \delta)})$ is a morphism in $\mathbf{SetF}(\Omega)$.

We define a morphism $u = u_{(A, \delta)} : (A, \delta) \rightarrow H(A, \delta) = (\text{w-singl}(A, \delta), \rho_{(A, \delta)})$ in $\mathbf{Set}(\Omega)$ such that $u(a, s) = F([-])(a, s)$ for all $a \in A, s \in \text{w-singl}(A, \delta)$, where F is a functor from Lemma 1.

Let $f : (A, \delta) \rightarrow (\text{w-singl}(B, \gamma), \rho_{(B, \gamma)})$ be a morphism in $\mathbf{Set}(\Omega)$. Then we define a morphism $\tilde{f} : (\text{w-singl}(A, \delta), \rho_{(A, \delta)}) \rightarrow (\text{w-singl}(B, \gamma), \rho_{(B, \gamma)})$ such that

$$\tilde{f}(s, t) = \bigvee_{x \in A} f(x, t) \otimes u(x, s),$$

for all $s \in \text{w-singl}(A, \delta), t \in \text{w-singl}(B, \gamma)$. Then \tilde{f} is a morphism in $\mathbf{Set}(\Omega)$. In fact, we have

$$\begin{aligned} \bigvee_{t \in \text{w-singl}(B, \gamma)} \tilde{f}(s, t) &= \bigvee_{x \in A} \left(\bigvee_{t \in \text{w-singl}(B, \gamma)} f(x, t) \right) \otimes u(x, s) = \\ &= \bigvee_{x \in A} u(x, s) = \bigvee_{x \in A} \bigwedge_{y \in A} \delta(x, y) \leftrightarrow s(y, 1) \geq \bigvee_{x \in A} s(x, 1) = 1. \end{aligned}$$

Moreover, the following diagram commutes.

$$\begin{array}{ccc} (A, \delta) & \xrightarrow{u} & (\text{w-singl}(A, \delta), \rho_{(A, \delta)}) \\ f \downarrow & & \downarrow \tilde{f} \\ (\text{w-singl}(B, \gamma), \rho_{(B, \gamma)}) & \xlongequal{\quad} & (\text{w-singl}(B, \gamma), \rho_{(B, \gamma)}) \end{array}$$

In fact, we have

$$\begin{aligned} \tilde{f} \circ u(a, t) &= \bigvee_{s \in \text{w-singl}(A, \delta)} u(a, s) \otimes \tilde{f}(s, t) = \\ &= \bigvee_{s \in \text{w-singl}(A, \delta)} u(a, s) \otimes \bigvee_{x \in A} u(x, s) \otimes f(x, t) \geq u(a, [a]) \otimes \bigvee_{x \in A} u(x, [a]) \otimes f(x, t) = \\ &= \bigvee_{x \in A} u(x, [a]) \otimes f(x, t) \geq f(a, t), \end{aligned}$$

and on the other hand

$$\tilde{f} \circ u(a, t) \leq \bigvee_{x \in A} \rho_{(A, \delta)}([a], [x]) \otimes f(x, t) = \bigvee_{x \in A} \delta(x, a) \otimes f(x, t) \leq f(a, t).$$

Let $f : (A, \delta) \rightarrow (B, \gamma)$ be a morphism in $\mathbf{Set}(\Omega)$. The morphism $H(f)$ will be defined as $\widetilde{u_{(B, \gamma)}} \circ f$.

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